Quotient Types by Idempotent Functions in Cedille

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Guide What is a Type Theory?

- Foundational theory for logic, math and computer science
- Commonly built on the lambda calculus
- Most common core is the Calculus of Constructions (CoC) (with function types, type quantification, etc)

Traditional Type Theory

(e.g. Agda and Coq)

CoCDependent Pair Inductive Identity Type $Id_A a b$ Inductive Definitions

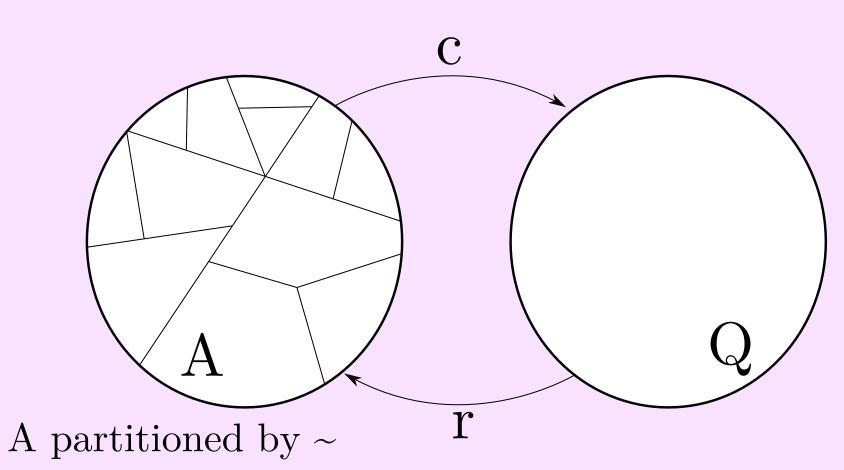
Cedille

Implicit (or erased) Function Types

U Dependent Intersections

Inductive Definitions are derived from these components

What is a Definable Quotient Type?



Given:

 $\forall a, b \in A. \text{if } a \sim b \text{ then } c(a) = c(b)$ $\forall a \in A.r(c(a)) \sim a$ $\forall q \in Q.c(r(q)) = q$

Then: Q is a quotient of A by ~

Given A and ~

Define: $f: A \to A$ such that $\forall a, b \in A. \text{if } a \sim b \text{ then } f(a) = f(b)$ $\forall a \in A.f(a) \sim a$

Then $Q = \Sigma a \in A.Id_A f(a) a$

Example:

 $(\mathbb{N},\mathbb{N}) \to (\mathbb{N},\mathbb{N})$ $f = \frac{\text{reduce to}}{}$ irreducible fraction

 $\mathbb{Q} = \Sigma(n,d) \in (\mathbb{N},\mathbb{N}).Id_{(\mathbb{N},\mathbb{N})} f((n,d)) (n,d)$

Define $f: A \to A$ Given A such that $\forall a \in A.f(f(a)) = f(a)$ Then $Q = \iota a \in A.f(a) \simeq a$

Equivalent to traditional definition

Key Difference: \sim is defined by fExample:

> $List\ A \rightarrow List\ A$ f = sort

 $SortedList\ A = \iota \ell \in List\ A.f(\ell) \simeq \ell$

Discussion of Definable Quotient Types

- Not all quotients are definable in the above sense (e.g. unordered pairs)
- Not all definable quotients can be defined as inductive types
- Quotients give a different perspective on constructing types

Quotients give two different meanings for a concept:

> Setoid: (\mathbb{N}, \mathbb{N}) and \sim Set: \mathbb{Q} and $Id_{\mathbb{Q}}$

Proofs can benefit from both

Projections can get in the way of equalities between views:

$$f((n,m)) = (n,m)$$
$$(n,m) = \pi_1 q$$

where π_1 is the first projection

All the same benefits and Q is a subtype of A

$$\forall q \in Q. \exists a \in A. q = a$$

This prevents projections from interfering with equalities between Q and A

$$f((n,m)) = (n,m)$$
$$(n,m) = q$$

Conclusions

- 1. Quotients by idempotent functions hide the equivalence relation ~
- 2. Cedille's encoding of quotients by idempotent functions gives you, additionally, subtypes

Related Work

- 1. Nuo Li, Quotient Types in Type Theory
- 2. Cyril Cohen, Pragmatic Quotient Types in Coq
- 3. Homotopy Type Theory for an alternative approach to quotients in a nontraditional theory